



Diagonal pivoting methods for solving tridiagonal systems without interchanges

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Summary

Problem: Solve

$$Tx = b,$$

where T is unsymmetric, tridiagonal, and nonsingular.

Challenge:

- Some applications (e.g., bi-conjugate gradient) do not allow interchanges
- GEPP: stable but requires interchanges

Approach: Use LBM^T where L and M are lower triangular and B is block diagonal with 1×1 and 2×2 blocks. Can demonstrate **backward stability**.



Background

Symmetric tridiagonal:

- Bunch (1974)
- Bunch and Kaufman (1977)
- Bunch and RM (2005, 2006)
- Fang and O'Leary (2006)

Composite-step BiCG:

- Bank and Chan (1994)

Stability analysis:

- Higham (lots)



Notation

Let

$$T = \begin{bmatrix} \alpha_1 & \gamma_2 & 0 & \cdots & 0 \\ \beta_2 & \alpha_2 & \gamma_3 & \ddots & \vdots \\ 0 & \beta_3 & \alpha_3 & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & \gamma_n \\ 0 & \cdots & 0 & \beta_n & \alpha_n \end{bmatrix}$$

and

$$\|T\|_{\max} = \max_{i,j} |T_{i,j}|.$$



Tridiagonal matrices

Partition

$$T = \begin{array}{cc} & \begin{array}{cc} d & n-d \end{array} \\ \begin{array}{c} d \\ n-d \end{array} & \left[\begin{array}{cc} B_1 & T_{12}^T \\ T_{21} & T_{22} \end{array} \right] \end{array}$$

T nonsingular $\Rightarrow B_1$ nonsingular for $d = 1$ or 2 : Note that

$$B_1 = [\alpha_1] \quad \text{or} \quad \begin{bmatrix} \alpha_1 & \gamma_2 \\ \beta_2 & \alpha_2 \end{bmatrix}$$

Then

$$\begin{aligned} \alpha_1 = 0 \quad \text{and} \quad \Delta = \alpha_1\alpha_2 - \beta_2\gamma_2 = 0 &\implies \beta_2 = 0 \quad \text{or} \quad \gamma_2 = 0 \\ &\implies T \text{ has a zero column or row} \\ &\implies T \text{ is singular.} \end{aligned}$$



Tridiagonal matrices

Partition

$$T = \begin{array}{c} d \\ n-d \end{array} \begin{array}{cc} d & n-d \\ \left[\begin{array}{cc} B_1 & T_{12}^T \\ T_{21} & T_{22} \end{array} \right] \end{array}$$

T nonsingular $\Rightarrow B_1$ nonsingular for $d = 1$ or 2 . Thus

$$T = \begin{bmatrix} I_d & 0 \\ T_{21}B_1^{-1} & I_{n-d} \end{bmatrix} \begin{bmatrix} B_1 & 0 \\ 0 & T_{22} - T_{21}B_1^{-1}T_{12}^T \end{bmatrix} \begin{bmatrix} I_d & B_1^{-1}T_{12}^T \\ 0 & I_{n-d} \end{bmatrix}.$$



Tridiagonal matrices

Partition

$$T = \begin{matrix} & d & n-d \\ & & \\ d & & \\ & & \\ n-d & & \end{matrix} \begin{bmatrix} B_1 & T_{12}^T \\ T_{21} & T_{22} \end{bmatrix}$$

T nonsingular $\Rightarrow B_1$ nonsingular for $d = 1$ or 2 . Thus

$$T = \begin{bmatrix} I_d & 0 \\ T_{21}B_1^{-1} & I_{n-d} \end{bmatrix} \begin{bmatrix} B_1 & 0 \\ 0 & T_{22} - T_{21}B_1^{-1}T_{12}^T \end{bmatrix} \begin{bmatrix} I_d & B_1^{-1}T_{12}^T \\ 0 & I_{n-d} \end{bmatrix}.$$

Schur complement $T_{22} - T_{21}B_1^{-1}T_{12}^T$ is tridiagonal:

$$T_{22} - \frac{\beta_2\gamma_2}{\alpha_1} e_1^{(n-1)} e_1^{(n-1)T} \quad \text{or} \quad T_{22} - \left(\frac{\alpha_1\beta_3\gamma_3}{\Delta} \right) e_1^{(n-2)} e_1^{(n-2)T}.$$



Tridiagonal matrices

Partition

$$T = \begin{matrix} & d & n-d \\ & & \\ d & & \\ & & \\ n-d & & \end{matrix} \begin{bmatrix} B_1 & T_{12}^T \\ T_{21} & T_{22} \end{bmatrix}$$

T nonsingular $\Rightarrow B_1$ nonsingular for $d = 1$ or 2 . Thus

$$T = \begin{bmatrix} I_d & 0 \\ T_{21}B_1^{-1} & I_{n-d} \end{bmatrix} \begin{bmatrix} B_1 & 0 \\ 0 & T_{22} - T_{21}B_1^{-1}T_{12}^T \end{bmatrix} \begin{bmatrix} I_d & B_1^{-1}T_{12}^T \\ 0 & I_{n-d} \end{bmatrix}.$$

Compute recursively to obtain

$$T = LBM^T$$

where L and M are lower triangular and B is block diagonal with 1×1 and 2×2 blocks. **Question:** How to choose size of pivot d ?



Proposed pivoting strategy

If $d = 1$, then

$$T = \left[\begin{array}{c|c} 1 & \\ \hline \beta_2/\alpha_1 & L_2 \end{array} \right] \left[\begin{array}{c|c} \alpha_1 & \\ \hline & B_2 \end{array} \right] \left[\begin{array}{c|c} 1 & \gamma_2/\alpha_1 \\ \hline & M_2^T \end{array} \right]$$

If $d = 2$, then

$$T = \left[\begin{array}{cc|c} 1 & & \\ & 1 & \\ \hline -\frac{\beta_2\beta_3}{\Delta} & \frac{\alpha_1\beta_3}{\Delta} & L_2 \end{array} \right] \left[\begin{array}{cc|c} \alpha_1 & \gamma_2 & \\ \beta_2 & \alpha_2 & \\ \hline & & B_2 \end{array} \right] \left[\begin{array}{cc|c} 1 & -\frac{\gamma_2\gamma_3}{\Delta} & \\ & 1 & \frac{\alpha_1\gamma_3}{\Delta} \\ \hline & & M_2^T \end{array} \right]$$

where $\Delta = \alpha_1\alpha_2 - \beta_2\gamma_2$.

Choose pivot size that yields smaller elements in L and M .



Proposed pivoting strategy

If $d = 1$, then

$$T = \left[\begin{array}{c|c} 1 & \\ \hline \beta_2/\alpha_1 & L_2 \end{array} \right] \left[\begin{array}{c|c} \alpha_1 & \\ \hline & B_2 \end{array} \right] \left[\begin{array}{c|c} 1 & \gamma_2/\alpha_1 \\ \hline & M_2^T \end{array} \right]$$

If $d = 2$, then

$$T = \left[\begin{array}{cc|c} 1 & & \\ & 1 & \\ \hline -\frac{\beta_2\beta_3}{\Delta} & \frac{\alpha_1\beta_3}{\Delta} & L_2 \end{array} \right] \left[\begin{array}{cc|c} \alpha_1 & \gamma_2 & \\ \beta_2 & \alpha_2 & \\ \hline & & B_2 \end{array} \right] \left[\begin{array}{cc|c} 1 & & -\frac{\gamma_2\gamma_3}{\Delta} \\ & 1 & \frac{\alpha_1\gamma_3}{\Delta} \\ \hline & & M_2^T \end{array} \right]$$

where $\Delta = \alpha_1\alpha_2 - \beta_2\gamma_2$. **Criterion #1:** Choose a 1×1 pivot if

$$\max \left\{ \frac{|\beta_2|}{|\alpha_1|}, \frac{|\gamma_2|}{|\alpha_1|} \right\} \leq \max \kappa \left\{ \frac{|\beta_2\beta_3|}{|\Delta|}, \frac{|\alpha_1\beta_3|}{|\Delta|}, \frac{|\gamma_2\gamma_3|}{|\Delta|}, \frac{|\alpha_1\gamma_3|}{|\Delta|} \right\}.$$



Proposed pivoting strategy

Criterion #1: Equivalently, choose a 1×1 pivot if

$$|\Delta| \cdot \max\{|\beta_2|, |\gamma_2|\} \leq |\alpha_1| \cdot \max\{\kappa\{|\beta_2\beta_3|, |\alpha_1\beta_3|, |\gamma_2\gamma_3|, |\alpha_1\gamma_3|\}\}$$

(**Intuition:** a 2×2 pivot is closer to being singular than a 1×1 pivot is to 0)

Criterion #2: Choose a 1×1 pivot if

$$|\alpha_1\alpha_2| \geq \kappa|\beta_2\gamma_2|.$$

(LBM^T reduces to the LDM^T factorization if T is positive definite).

Constant: $\kappa = (\sqrt{5} - 1)/2 \approx 0.62$ (Bunch (1974))



Proposed pivoting strategy

Algorithm 1.

$$\kappa = (\sqrt{5} - 1)/2 \approx 0.62$$

$$\Delta = \alpha_1\alpha_2 - \beta_2\gamma_2$$

if $|\Delta| \cdot \max\{|\beta_2|, |\gamma_2|\} \leq |\alpha_1| \cdot \max\{\kappa\{|\beta_2\beta_3|, |\alpha_1\beta_3|, |\gamma_2\gamma_3|, |\alpha_1\gamma_3|\}\}$
or $|\alpha_1\alpha_2| \geq \kappa|\beta_2\gamma_2|$

$$d_I = 1$$

else

$$d_I = 2$$

end



Alternative pivoting strategy

Algorithm 2.

$$\kappa = (\sqrt{5} - 1)/2 \approx 0.62$$

$$\sigma_1 = \max\{|\alpha_2|, |\gamma_2|, |\beta_2|, |\gamma_3|, |\beta_3|\}$$

if $|\alpha_1|\sigma_1 \geq \kappa|\beta_2\gamma_2|$,

$$d_{II} = 1$$

else

$$d_{II} = 2$$

end

$$T = \begin{bmatrix} \alpha_1 & \gamma_2 & 0 & \cdots & 0 \\ \beta_2 & \alpha_2 & \gamma_3 & \ddots & \vdots \\ 0 & \beta_3 & \alpha_3 & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & \gamma_n \\ 0 & \cdots & 0 & \beta_n & \alpha_n \end{bmatrix}$$

(cf. Bunch-Kaufman strategy for symmetric tridiagonal matrices – different from the more general Bunch-Kaufman LBL^T).



Pivoting strategies

Lemma. *If $d_I = 1$, then $d_{II} = 1$. If $d_{II} = 2$, then $d_I = 2$. And, if $d_I = 2$ and $d_{II} = 1$, then the subsequent pivot size \tilde{d}_{II} is 1.*

Summary: The two pivoting strategies differ only when the first strategy chooses a 2×2 pivot while the alternative chooses two 1×1 pivots.



Main theorem

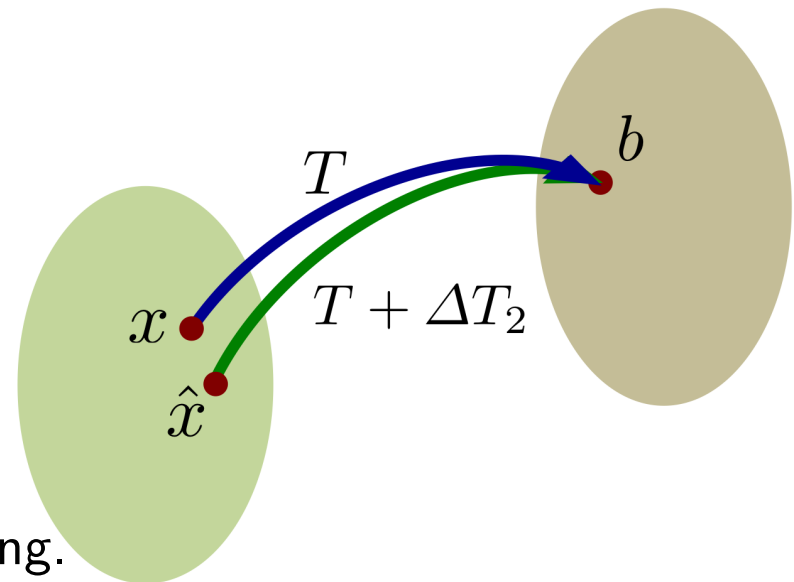
Theorem. Let $T \approx \hat{L}\hat{B}\hat{M}^T$ be the computed LBM^T factorization of the unsymmetric tridiagonal matrix $T \in \mathbb{R}^{n \times n}$ using Algorithm 1 or 2. Assume that all linear systems $Ey = f$ involving 2×2 pivots E are solved using the explicit inverse. Then

$$T - \hat{L}\hat{B}\hat{M}^T = \Delta T_1, \quad \text{and} \quad (T + \Delta T_2)\hat{x} = b,$$

where $\|\Delta T_i\|_{\max} \leq cu\|T\|_{\max} + O(u^2)$,
for $i = 1, 2$, where c is a constant.

Summary: \hat{x} is the exact solution to a nearby problem.

Proof: Long and not particularly enlightening.





Numerical results

Obtained 16 unsymmetric tridiagonal matrices from Dhillon (1998) and Hargreaves (2004) on estimating condition numbers of tridiagonal matrices.

Matrix type	Description
1	Random elements (uniformly distributed on $[-1, 1]$).
2	Preassigned condition number with one very small singular value: <code>gallery('randsvd',n,1e15,2,1,1)</code> .
3	Preassigned condition number with geometric distribution of singular values: <code>gallery('randsvd',n,1e15,3,1,1)</code> .
7	Ill-conditioned, tridiagonal: <code>gallery('dorr',n,1e-4)</code> .
12	Toeplitz: main diagonal zero, each off-diagonal a random number (uniformly distributed on $[-1, 1]$).
15	Tridiagonal with main diagonal zero and known eigenvalues: <code>gallery('clement',n,0)</code> .



Numerical results

Generated 50×50 matrices of each type and computed $\|A\hat{x} - b\|_2$:

Matrix Type	Algorithm 1	Algorithm 2	GEPP	Cond. Num.
1	1.437796e-014	1.289799e-014	8.567857e-015	2.942402e+002
2	5.012055e-003	5.012055e-003	1.121297e-003	8.731214e+014
3	1.380824e-003	1.108818e-003	8.331566e-004	1.001503e+015
7	3.328680e+002	3.328680e+002	5.128745e+002	7.574961e+016
12	4.198353e-010	4.198353e-010	1.810847e-010	5.673349e+007
15	1.753491e-001	1.753491e-001	1.004420e-001	3.449791e+015

Conclusion: GEPP performs slightly better than LBM^T , but results are comparable.



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Approach: Use LBM^T where L and M are lower triangular and B is block diagonal with 1×1 and 2×2 blocks. Can demonstrate **backward stability**.

Future work: Extension to unsymmetric Lanczos / bi-conjugate gradient method.



References

A backward stability analysis of diagonal pivoting methods for solving unsymmetric tridiagonal systems without interchanges, Jennifer Erway and RM, Accepted for publication in Numerical Linear Algebra with Applications.

Solving unsymmetric tridiagonal systems without interchanges, Joseph Tyson, Jennifer Erway, and RM, In preparation.

google: roummel



Thank you.