# EECS 275 Matrix Computation 

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Lecture 15

## Overview

- Inverse iteration
- Rayleigh quotient iteration
- Condition
- Perturbation
- Stability


## Reading

- Chapter 27 and Chapter 12 of Numerical Linear Algebra by Llyod Trefethen and David Bau
- Chapter 7 of Matrix Computations by Gene Golub and Charles Van Loan


## Inverse iteration

- For any $\mu \in \mathbb{R}$ that is not an eigenvalue of $A$, the eigenvectors of $(A-\mu I)^{-1}$ are the same as the eigenvectors of $A$, and the the corresponding eigenvalues are $\left(\lambda_{j}-\mu\right)^{-1}$ where $\lambda_{j}$ are the eigenvalues of $A$
- Suppose $\mu$ is close to an eigenvalue $\lambda_{J}$ of $A$, then $\left(\lambda_{J}-\mu\right)^{-1}$ may be much larger than $\left(\lambda_{j}-\mu\right)^{-1}$ for all $j \neq J$
- If we apply power iteration to $(A-\mu I)^{-1}$, the process will converge rapidly to $\mathbf{q}_{J}$
- Algorithm:

Initialize $\mathbf{v}^{(0)}$ randomly with $\left\|\mathbf{v}^{(0)}\right\|=1$
for $k=1,2, \ldots$ do
Solve $\mathbf{w}=(A-\mu I)^{-1} \mathbf{v}^{(k-1)} \quad / /$ apply $(A-\mu I)^{-1}$
$\mathbf{v}^{(k)}=\frac{\mathbf{w}}{\|\mathbf{w}\|} \quad / /$ normalize
$\lambda^{(k)}=\left(\mathbf{v}^{(k)}\right)^{\top} A \mathbf{v}^{(k)} \quad / /$ Rayleigh quotient
end for

## Inverse iteration (cont'd)

- what if $\mu$ is or nearly is an eigenvalue of $A$, so that $A-\mu /$ is singular?
- can be handled numerically
- exhibit linear convergence
- unlike power iteration, we can choose the eigenvector that will be found by supplying an estimate $\mu$ of the corresponding eigenvalue
- if $\mu$ is much closer to one eigenvalue of $A$ than to the others, then the largest eigenvalue o $(A-\mu I)^{-1}$ will be much larger than the rest


## Theorem

Suppose $\lambda_{J}$ is the closest eigenvalue of $\mu$ and $\lambda_{K}$ is the second closest, that is $\left|\mu-\lambda_{J}\right|<\left|\mu-\lambda_{K}\right|<\left|\mu-\lambda_{j}\right|$ for each $j \neq J$. Furthermore, suppose $\mathbf{q}_{J}^{\top} \mathbf{v}^{(0)} \neq 0$, then after $k$ iterations

$$
\left\|\mathbf{v}^{(k)}-\left( \pm \mathbf{q}_{J}\right)\right\|=O\left(\left|\frac{\mu-\lambda_{J}}{\mu-\lambda_{K}}\right|^{k}\right),\left|\lambda^{(k)}-\lambda_{J}\right|=O\left(\left|\frac{\mu-\lambda_{J}}{\mu-\lambda_{K}}\right|^{2 k}\right)
$$

as $k \rightarrow \infty$. The $\pm$ sign means that at each step $k$, one or the other choice of sign is to be taken, and then the indicated bound holds

## Inverse iteration and Rayleigh quotient iteration

- Inverse iteration:
- one of the most valuable tools of numerical linear algebra
- a standard method of calculating one or more eigenvectors of a matrix if the eigenvalues are already known
- Obtaining an eigenvalue estimate from an eigenvector estimate (the Rayleigh quotient)
- Obtaining an eigenvector estimate from an eigenvalue estimate (inverse iteration)
- Rayleigh quotient algorithm:

Initialize $\mathbf{v}^{(0)}$ randomly with $\left\|\mathbf{v}^{(0)}\right\|=1$ $\lambda^{(0)}=\left(\mathbf{v}^{(0)}\right)^{\top} A \mathbf{v}^{(0)}=$ corresponding Rayleigh quotient
for $k=1,2, \ldots$ do
Solve $\mathbf{w}=\left(A-\lambda^{(k-1)} /\right)^{-1} \mathbf{v}^{(k-1)} \quad / /$ apply $\left(A-\mu^{(k-1)} I\right)^{-1}$ $\mathbf{v}^{(k)}=\frac{\mathbf{w}}{\|\mathbf{w}\|} \quad / /$ normalize
$\lambda^{(k)}=\left(\mathbf{v}^{(k)}\right)^{\top} A \mathbf{v}^{(k)} \quad / /$ Rayleigh quotient
end for

## Rayleigh quotient iteration

## Theorem

Rayleigh quotient iteration converges to an eigenvalue/eigenvector pair for all except a set of measure zero of starting vectors $\mathbf{v}^{(0)}$. When it converges, the convergence is ultimately cubic in the sense that if $\lambda_{J}$ is an eigenvalue of $A$ and $\mathbf{v}^{(0)}$ is sufficiently close to the eigenvector $\mathbf{q}_{J}$, then

$$
\left\|\mathbf{v}^{(k+1)}-\left( \pm \mathbf{q}_{J}\right)\right\|=O\left(\left\|\mathbf{v}^{(k)}-\left( \pm \mathbf{q}_{J}\right)\right\|^{3}\right)
$$

and

$$
\left|\lambda^{(k+1)}-\lambda_{J}\right|=O\left(\left|\lambda^{(k)}-\lambda_{J}\right|^{3}\right)
$$

as $k \rightarrow \infty$. The $\pm$ sign means that at each step $k$, one or the other choice of sign is to be taken, and then the indicated bound holds.

## Example

- Consider the symmetric matrix

$$
A=\left[\begin{array}{lll}
2 & 1 & 1 \\
1 & 3 & 1 \\
1 & 1 & 4
\end{array}\right]
$$

and let $\mathbf{v}^{(0)}=(1,1,1)^{\top} / \sqrt{3}$ the initial eigenvector estimate

- When Rayleigh quotient iteration is applied to $A$, the following values $\lambda^{(k)}$ are computed by the first 3 iterations:

$$
\lambda^{(0)}=5, \quad \lambda^{(1)}=5.2131 \ldots, \quad \lambda^{(2)}=5.214319743184 \ldots
$$

The actual value of the eigenvalue corresponding to the eigenvector closest to $\mathbf{v}^{(0)}$ is $\lambda=5.214319743377$

- After three iterations, Rayleigh quotient iteration has produced a result accurate to 10 digits
- With three more iterations, the solution increases this figure to about 270 digits, if our machine precision were high enough


## Operation counts

- Flops (floating point operations): addition, subtraction, multiplication, division, or square root counts as one flop
- Suppose $A \in \mathbb{R}^{m \times m}$ is a full matrix, each step of power iteration involves a matrix-vector multiplication, requiring $O\left(m^{2}\right)$ flops
- Each step of inverse iteration involves the solution of a linear system, which might seem to require $O\left(m^{3}\right)$ flops, but this can be reduced to $O\left(m^{2}\right)$ if the matrix is processed in advance by LU or QR factorization
- In Rayleigh quotient iteration, the matrix to be inverted changes at each step, and beating $O\left(m^{3}\right)$ flops per step
- These figures improve greatly if $A$ is tridiagonal, and all three iterations requires just $O(m)$ flops per step
- For non-symmetric matrices, we must deal with Hessenberg instead of tridiagonal structure, and this figure increases to $O\left(\mathrm{~m}^{3}\right)$


## Conditioning and condition numbers

- Conditioning: pertain to the perturbation behavior of a mathematical problem
- Stability: pertain to the perturbation behavior of an algorithm used to solve the problem on a computer
- We view a problem as a function $f: X \rightarrow Y$ from a normed vector space $X$ of data to a normed vector space $Y$ of solutions
- $f$ is usually nonlinear (even in linear algebra), but most of the time it is at least continuous
- A well-conditioned problem: all small perturbations of $\mathbf{x}$ lead to only small changes in $f(\mathbf{x})$
- An ill-conditioned problem: some small perturbation of $\mathbf{x}$ leads to a large change in $f(\mathbf{x})$


## Absolute condition number

- Let $\delta \mathbf{x}$ denote a small perturbation of $\mathbf{x}$ and write $\delta f=f(\mathbf{x}+\delta \mathbf{x})-f(\mathbf{x})$
- Absolute condition number:

$$
\hat{\kappa}=\hat{\kappa}(\mathbf{x})=\lim _{\delta \rightarrow 0} \sup _{\|\delta \mathbf{x}\| \leq \delta} \frac{\|\delta f\|}{\|\delta \mathbf{x}\|}
$$

- For most problems, it can be interpreted as a supremum over all infinitesimal perturbations $\delta \mathbf{x}$, and can write simply as

$$
\hat{\kappa}=\sup _{\delta \mathbf{x}} \frac{\|\delta f\|}{\|\delta \mathbf{x}\|}
$$

- If $f$ is differentiable, we can evaluate the condition number by means of the derivative of $f$
- Let $J(\mathbf{x})$ be the matrix where $i, j$ entry is the partial derivative $\partial f_{i} / \partial \mathbf{x}_{j}$ evaluated at $\mathbf{x}$, known as the Jacobian of $f$ at $\mathbf{x}$
- By the definition, $\delta f \approx J(\mathbf{x}) \delta \mathbf{x}$ with $\|\delta \mathbf{x}\| \rightarrow 0$, and the absolute condition number becomes

$$
\hat{\kappa}=\|J(\mathbf{x})\|
$$

where $\|J(\mathbf{x})\|$ represents the norm of $J(\mathbf{x})$

## Relative condition number

- Relative condition number:

$$
\kappa=\lim _{\delta \rightarrow 0} \sup _{\|\delta \mathbf{x}\| \leq \delta}\left(\frac{\|\delta f\|}{\|f(\mathbf{x})\|} / \frac{\|\delta \mathbf{x}\|}{\|\mathbf{x}\|}\right)
$$

- Again, assume $\delta \mathbf{x}$ and $\delta f$ are infinitesimal,

$$
\kappa=\sup _{\delta \mathbf{x}}\left(\frac{\|\delta f\|}{\|f(\mathbf{x})\|} / \frac{\|\delta \mathbf{x}\|}{\|\mathbf{x}\|}\right)
$$

- If $f$ is differentiable, we can express this in terms of Jacobian

$$
\kappa=\frac{\|J(\mathbf{x})\|}{\|f(\mathbf{x})\| /\|\mathbf{x}\|}
$$

- Relative condition number is more important in numerical analysis
- A problem is well-conditioned if $\kappa$ is small (e.g., $1,10,10^{2}$ ), and ill-conditioned if $\kappa$ is large (e.g., $10^{6}, 10^{16}$ )


## Well-conditioned and ill-conditioned problems

- Consider the problem of obtaining the scalar $x / 2$ from $x \in \mathbb{C}$. The Jacobian of the function $f: x \rightarrow x / 2$ is the derivative $J=f^{\prime}=1 / 2$

$$
\kappa=\frac{\|J\|}{\|f(x)\| /\|x\|}=\frac{1 / 2}{(x / 2) / x}=1
$$

This problem is well-conditioned

- Consider $x^{2}-2 x+1=(x-1)^{2}$, with a double root $x=1$. A small perturbation in the coefficient may lead to large change in the roots

$$
x^{2}-2 x+0.9999=(x-0.99)(x-1.01)
$$

In fact, the roots can change in proportion to the square root of the change in the coefficients, so in this case the Jacobian is infinite (the problem is not differentiable), and $\kappa=\infty$

## Computing eigenvalues of a non-symmetric matrix

- Computing the eigenvalues of a non-symmetric matrix is often ill-conditioned
- Consider the two matrices

$$
A=\left[\begin{array}{cc}
1 & 1000 \\
0 & 1
\end{array}\right] \quad B=\left[\begin{array}{cc}
1 & 1000 \\
0.001 & 1
\end{array}\right]
$$

and the eigenvalues are $\{1,1\}$ and $\{0,2\}$ respectively

- On the other hand, if a matrix $A$ is symmetric (more generally, if it is normal), then its eigenvalues are well-conditioned
- It can be shown that if $\lambda$ and $\lambda+\delta \lambda$ are corresponding eigenvalues of $A$ and $A+\delta A$, then $|\delta \lambda| \leq\|\delta A\|_{2}$ with equality if $\delta A$ is a multiple of the identity
- In that case, $\hat{\kappa}=1$ if perturbations are measured in the 2-norm and the relative condition number is $\kappa=\|A\|_{2} /|\lambda|$


## Condition of matrix-vector multiplication

- Consider $A \in \mathbb{C}^{m \times n}$, determine a condition number corresponding to perturbations of $\mathbf{x}$

$$
\kappa=\sum_{\delta \mathbf{x}}\left(\frac{\|A(\mathbf{x}+\delta x)-A \mathbf{x}\|}{\|A \mathbf{x}\|} / \frac{\|\delta \mathbf{x}\|}{\|\mathbf{x}\|}\right)=\sup _{\delta \mathbf{x}} \frac{\|A \delta \mathbf{x}\|}{\|\delta \mathbf{x}\|} / \frac{\|A \mathbf{x}\|}{\|\mathbf{x}\|},
$$

that is

$$
\kappa=\|A\| \frac{\|\mathbf{x}\|}{\|A \mathbf{x}\|} \quad\left(\kappa=\frac{\|J(\mathbf{x})\|}{\|f(\mathbf{x})\| /\|\mathbf{x}\|}\right)
$$

which is an exact formula for $\kappa$, dependent on both $A$ and $\mathbf{x}$

- If $A$ happens to be square and non-singular, we can use the fact that $\|\mathbf{x}\| /\|A \mathbf{x}\| \leq\left\|A^{-1}\right\|$ to loosen the bound

$$
\kappa \leq\|A\|\left\|A^{-1}\right\|
$$

or

$$
\kappa=\alpha\|A\|\left\|A^{-1}\right\|, \quad \alpha=\frac{\|\mathbf{x}\|}{\|A \mathbf{x}\|} /\left\|A^{-1}\right\|
$$

## Condition of matrix-vector multiplication (cont'd)

- For certain choices of $\mathbf{x}$, we have $\alpha=1$, and consequently $\kappa=\|A\|\left\|A^{-1}\right\|$
- If we use 2-norm, this will occur when $\mathbf{x}$ is a multiple of a minimal right singular vector of $A$
- If $A \in \mathbb{C}^{m \times n}$ with $m \geq n$ has full rank, the above equations hold with $A^{-1}$ replaced by the pseudo-inverse $A^{\dagger}=\left(A^{H} A\right)^{-1} A^{H} \in \mathbb{C}^{n \times m}$
- Given $A$, compute $A^{-1} \mathbf{b}$ from input $\mathbf{b}$ ?
- Mathematically, identical to the problem just considered except that $A$ is repalced by $A^{-1}$


## Condition of matrix-vector multiplication (cont'd)

## Theorem

Let $A \in \mathbb{C}^{m \times n}$ be nonsingular and consider the equation $A \mathbf{x}=\mathbf{b}$. The problem of computing $\mathbf{b}$ given $\mathbf{x}$, has condition number

$$
\begin{equation*}
\kappa=\|A\| \frac{\|\mathbf{x}\|}{\|\mathbf{b}\|} \leq\|A\|\left\|A^{-1}\right\| \tag{1}
\end{equation*}
$$

with respect to perturbation of $\mathbf{x}$. The problem of computing $\mathbf{x}$ given $\mathbf{b}$, has condition number

$$
\begin{equation*}
\kappa=\left\|A^{-1}\right\| \frac{\|\mathbf{b}\|}{\|\mathbf{x}\|} \leq\|A\|\left\|A^{-1}\right\| \tag{2}
\end{equation*}
$$

with respect to perturbation of $\mathbf{b}$. If we use 2-norm, then the equality holds in (1) if $\mathbf{x}$ is a multiple of a right singular vector of $A$ corresponding to the minimal singular value $\sigma_{m}$. Likewise, the equality hold in (2) if $\mathbf{b}$ is a multiple of a left singular vector of A corresponding to the maximal singular value $\sigma_{1}$.

## Condition number of a matrix

- The product $\|A\|\left\|A^{-1}\right\|$ is the condition number of $A$ (relative to the norm $\|\cdot\|)$, denoted by $\kappa(A)$

$$
\kappa(A)=\|A\|\left\|A^{-1}\right\|
$$

- The condition number is attached to a matrix, not a problem
- If $\kappa(A)$ is small, $A$ is said to be well-conditioned; if $\kappa(A)$ is large, $A$ is ill-conditioned.
- If $A$ is singular, it is customary to write $\kappa(A)=\infty$
- If $\|\cdot\|=\|\cdot\|_{2}$, then $\|A\|=\sigma_{1}$, and $\left\|A^{-1}\right\|=1 / \sigma_{m}$, and thus

$$
\kappa(A)=\frac{\sigma_{1}}{\sigma_{m}}
$$

in the 2-norm, which is the formula for computing 2-norm condition numbers of matrices.

- The ratio $\sigma_{1} / \sigma_{m}$ can be interpreted as the eccentricity of the hyperellipse that is the image of the unit sphere of $\mathbb{C}^{m}$ under $A$


## Condition number of a matrix (cont'd)



- For a rectangular matrix $A \in C^{m \times n}$ of full rank, $m \geq n$, the condition number is defined in terms of pseudo-inverse

$$
\kappa(A)=\|A\|\left\|A^{\dagger}\right\|
$$

since $A^{\dagger}$ is motivated by least squares problem, this definition is most useful with 2-norm, where we have

$$
\kappa(A)=\frac{\sigma_{1}}{\sigma_{n}}
$$

## Condition number of a system of equations

- Let $\mathbf{b}$ be fixed, and consider the behavior of the problem $A \mapsto \mathbf{x}=A^{-1} \mathbf{b}$ when $A$ is perturbed by infinitesimal $\delta A$, then

$$
(A+\delta A)(\mathbf{x}+\delta \mathbf{x})=\mathbf{b}
$$

- Using $A \mathbf{x}=\mathbf{b}$ and dropping the term $(\delta A)(\delta \mathbf{x})$, we obtain $(\delta A) \mathbf{x}+A(\delta \mathbf{x})=0$, that is, $\delta \mathbf{x}=-A^{-1}(\delta A) \mathbf{x}$
- It implies $\|\delta \mathbf{x}\| \leq\left\|A^{-1}\right\|\|\delta A\|\|\mathbf{x}\|$, or equivalently

$$
\frac{\|\delta \mathbf{x}\|}{\|\mathbf{x}\|} / \frac{\|\delta A\|}{\|A\|} \leq\left\|A^{-1}\right\|\|A\|=\kappa(A)
$$

- The equality in this bound will hold when $\delta A$ is such that

$$
\left\|A^{-1}(\delta A) \mathbf{x}\right\|=\left\|A^{-1}\right\|\|\delta A\|\|\mathbf{x}\|
$$

and it can be shown that for any $A$ and $\mathbf{b}$ and norm $\|\cdot\|$, such perturbations of $\delta A$ exists

## Condition number of a system of equations (cont'd)

## Theorem

Let $\mathbf{b}$ be fixed and consider the problem of computing $\mathbf{x}=A^{-1} \mathbf{b}$ where $A$ is square and nonsingular. The condition number of this problem with respect to perturbation in $A$ is

$$
\kappa=\|A\|\left\|A^{-1}\right\|=\kappa(A)
$$

- Both theorems regarding matrix-vector multiplication and system of equations are of fundamental importance in numerical linear algebra
- They determine how accurately one can solve systems of equations
- If a problem $A \mathbf{x}=\mathbf{b}$ contains an ill-conditioned matrix $A$, one must always expect to "lose $\log _{10} \kappa(A)$ digits" in computing the solution (except under very special cases)

