Total possible marks: 100. Homeworks must be solved individually. This set covers chapters 6–12 of the textbook Introduction to Algorithms, 3rd. ed., by Cormen et al.

Exercise 1: heapsort (26 points). We consider Heapsort to sort an array in decreasing order by using min-heaps.

1. (3 points) State the min-heap property.

Then, using the same notation as in the textbook, write pseudocode for the following functions (where A is the min-heap):

2. (3 points) Min-Heapify(A, i), which assumes that the binary trees rooted at \( \text{LEFT}(i) \) and \( \text{RIGHT}(i) \) are min-heaps, but that \( A[i] \) may be larger than its children. Min-Heapify lets the value of \( A[i] \) float down so that the subtree rooted at \( i \) obeys the min-heap property.

3. (2 points) Build-Min-Heap(A), which builds a min-heap on the input array \( A \) (overwriting it).

4. (2 points) Heapsort(A), which sorts the array \( A \) in decreasing order, based on Min-Heapify and Build-Min-Heap.

Finally:

5. (10 points) State a loop invariant for Heapsort and use it to prove its correctness. Assume Min-Heapify and Build-Min-Heap are correct.

6. (6 points) Give the runtime for Min-Heapify, Build-Min-Heap and Heapsort as a function of the array size \( n \). Explain your answers.

Exercise 2: counting sort (15 points). Write the pseudocode of an algorithm that, given an array \( A \) of \( n \) integers in \( \{0, \ldots, k\} \), preprocesses \( A \) in \( \Theta(n + k) \) time and then answers any query about how many integers fall into the range \( [a, b] \) in \( \Theta(1) \) time.

Exercise 3: hash tables (27 points). Consider a hash table with \( m = 11 \) slots and using the hash function \( h(k) = \lfloor m(kA - \lfloor kA \rfloor) \rfloor \) where \( A = (\sqrt{5} - 1)/2 \) and \( k \) is a natural number. Consider the keys \( k = 10, 22, 31, 4, 15, 28, 17, 88, 59 \) (in that order).

1. (3 points) Give \( h(k) \) for each of those keys.

Now, consider inserting those keys in the order given above into the hash table. Show the final table in these two cases:

2. (12 points) Chaining using as hash function \( h(k) \).

3. (12 points) Open addressing using linear probing and the same hash function \( h(k) \).
Exercise 4: hash tables (10 points). We have a hash table $T_1$ that uses chaining to resolve collisions; it has $m$ slots and contains currently $n$ elements. We have another hash table $T_2$ that uses open addressing; it contains currently $n$ elements and occupies exactly the same amount of memory as $T_1$. Now, we execute a search for a key $k$ which is in neither table. On average, in which of the two tables does the search take fewer operations? Hint: assume that pointers and keys occupy each one word of memory, and that an unsuccessful search requires on average $1 + \alpha(T_1)$ operations for $T_1$ and $\frac{1}{1-\alpha(T_2)}$ operations for $T_2$, where $\alpha$ is the load factor.

Exercise 5: binary search trees (22 points).

1. (10 points) Starting from an empty binary search tree, draw the final tree resulting from the insertion of the following keys: 16, 25, 4, 12, 10, 5, 15, 42, 66, 2 (in that order).

2. (2 points) Do the inorder tree walk, printing the resulting keys.

3. (10 points) Starting from the tree obtained in the former question, draw the tree resulting from the removal of the following keys: 4, 12 (in that order).

Bonus exercise: bucket sort (20 points). (Exercise 8.4-4 in the book.) We are given $n$ points in the unit circle, $p_i = (x_i, y_i)$, such that $0 < x_i^2 + y_i^2 \leq 1$ for $i = 1, 2, \ldots, n$. Suppose that the points are uniformly distributed; that is, the probability of finding a point in any region of the circle is proportional to the area of that region. Design an algorithm with an average-case running time of $\Theta(n)$ to sort the $n$ points by their distances $d_i = \sqrt{x_i^2 + y_i^2}$ from the origin.

Hint: Design the bucket sizes in BUCKET-SORT to reflect the uniform distribution of the points in the unit circle.